

Verified Bounds for the p -Norm Condition Number*

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Abstract

Methods to compute verified error bounds for the p -norm condition number of a matrix are discussed for $p \in \{1, 2, \infty\}$ and the Frobenius norm. We consider the cases of a real or complex, point or interval input matrix. In the latter case the condition number of all matrices within the interval matrix are bounded. A special method for extremely ill-conditioned matrices is derived as well. Numerical results suggest that the quality of the bounds corresponds to the fact that the condition number of the condition number is the condition number.

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1 Introduction and Notation

Let a real or complex square matrix A be given, that is $A \in \mathbb{K}^{n \times n}$ where $\mathbb{K} \in \{\mathbb{R}, \mathbb{C}\}$. When using $\|\cdot\|_p$ we henceforth assume $p \in \{1, 2, \infty, F\}$, where $\|A\|_F$ denotes the Frobenius norm $\sqrt{\sum a_{ij}^2}$. We derive methods to compute verified error bounds for $\kappa_p(A) := \|A\|_p \|A^{-1}\|_p$. Note that finite bounds for the condition number imply A to be non-singular. We also consider the case of an interval matrix $\mathbf{A} \in \mathbb{IK}^{n \times n}$, where all condition numbers of $A \in \mathbf{A}$ are bounded. Moreover, a special method for extremely ill-conditioned matrices is presented. We assume the reader to be familiar with basic concepts of interval arithmetic, cf. [5].

A straightforward INTLAB [7] statement to compute an inclusion $\mathbf{Cnd} \in \mathbb{IR}$ of $\kappa_p(A)$ is

$$\mathbf{Cnd} = \mathbf{norm}(\mathbf{intval}(\mathbf{A}), p) * \mathbf{norm}(\mathbf{inv}(\mathbf{intval}(\mathbf{A})), p) \quad (1)$$

which works both for point and interval, real and complex matrices \mathbf{A} . It fails for too ill-conditioned matrices with $\mathbf{Cnd} = \text{NaN}$. In the following we discuss alternative and

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specialized methods to compute an inclusion of the condition number. We frequently need

$$A, B \in \mathbb{K}^{n \times n} \text{ and } |A| \leq B \quad \Rightarrow \quad \|A\|_p \leq \|B\|_p, \quad (2)$$

where comparison is understood to be componentwise. This is clear for $p \in \{1, \infty, F\}$. For $p = 2$ use $\|A\|_2^2 \leq \| |A| \|_2^2 = \rho(|A|^T |A|) \leq \rho(B^T B) = \|B\|_2^2$, where $\rho(\cdot)$ denotes the spectral radius and the latter inequality follows by Perron-Frobenius Theory [4].

2 Main Results

Bounds for the p -norm of a matrix with $p \in \{1, \infty, F\}$ are immediate using interval arithmetic, and bounds for the spectral norm are obtained by the methods in [8]. So we concentrate on bounds for the norm of the inverse of a matrix A .

Such bounds are computed on the basis of an *approximate* inverse R of A . The quality of the bounds for $\|A^{-1}\|$ relates directly to the quality of the approximation $R \approx A^{-1}$. This quality depends, in turn, on the condition number of the matrix A , which is to be bounded and which is not known a priori.

What sounds like a vicious circle is resolved by computing a bound α for the norm of $E := I - RA$. If $\alpha < 1$, then lower and upper bounds for the norm of the true inverse A^{-1} follow using the norm of R . The norm of R is bounded using interval arithmetic as we explained previously for the norm of A .

Note that mathematically $R \in \mathbb{F}^{n \times n}$ is an arbitrary matrix. If R is of poor quality, the condition $\alpha < 1$ is not satisfied. In practice R is computed by some standard numerical algorithm, usually based on Gaussian elimination with partial pivoting. Numerically, we can expect $\alpha < 1$ if $\kappa(A) \lesssim \text{eps}^{-1}$ for eps denoting the relative rounding error unit (about $1.1 \cdot 10^{-16}$ in IEEE 754 binary64). It is important to note that all our bounds are valid without any *a priori* assumption on the condition number of A or the the quality of R .

2.1 Perturbations of the Identity Matrix

Denoting the identity matrix of proper dimension by I , we first investigate $\|(I - E)^{-1}\|_p$ for admissible perturbations E .

Lemma 2.1 *Let $E \in \mathbb{K}^{n \times n}$ and $p \in \{1, 2, \infty\}$ be given, and suppose $\|E\|_p \leq \alpha < 1$. Then*

$$\frac{1}{1 + \alpha} \leq \|(I - E)^{-1}\|_p \leq \frac{1}{1 - \alpha}. \quad (3)$$

Equality is attained in both bounds for $E = -\alpha I$ and $E = \alpha I$, respectively, for all p and any n .

PROOF. The result follows by using the Neumann series and

$$1 = \|(I - E)(I - E)^{-1}\|_p \leq (1 + \alpha)\|(I - E)^{-1}\|_p,$$

a commonly used “splitting of 1”. □

Note the apparent asymmetry of the bounds: For nonzero α , the lower bound is closer to 1 than the upper bound. This will be utilized later. For the Frobenius norm, things are a little more involved.

Lemma 2.2 *Let $E \in \mathbb{K}^{n \times n}$ be given, and suppose $\|E\|_F \leq \alpha < 1$. Then*

$$\frac{\sqrt{n}}{1 + \|E\|_2} \leq \|(I - E)^{-1}\|_F \leq \frac{\sqrt{n}}{1 - \|E\|_2} . \quad (4)$$

Equality is attained in both bounds for $\beta = n^{-1/2}\alpha$ and $E = -\beta I$ and $E = \beta I$, respectively, for any n . Moreover,

$$\sqrt{n} - \frac{\alpha}{1 - \alpha} \leq \|(I - E)^{-1}\|_F \leq \sqrt{n} + \frac{\alpha}{1 - \alpha} . \quad (5)$$

The left and right bound in (5) can be better up to a factor 2 and \sqrt{n} than the corresponding bounds in (4), respectively. Conversely, both bounds in (5) can be arbitrarily weaker than the bounds in (4).

PROOF. The standard perturbation bound for singular values (Theorem 3.3.16 in [3]) states that

$$A, E \in \mathbb{K}^{m \times n} \quad \Rightarrow \quad |\sigma_i(A + E) - \sigma_i(A)| \leq \|E\|_2 \quad \text{for } 1 \leq i \leq \min(m, n) .$$

Setting $m = n$ and $A := I$, it follows that for all $1 \leq i \leq n$

$$1 - \|E\|_2 \leq \sigma_i(I - E) \leq 1 + \|E\|_2 .$$

Then $\|E\|_2 \leq \|E\|_F$ and $\|E\|_F \leq \alpha < 1$ shows that $I - E$ has full rank and

$$\frac{1}{1 + \|E\|_2} \leq \sigma_i((I - E)^{-1}) \leq \frac{1}{1 - \|E\|_2} . \quad (6)$$

For any matrix $M \in \mathbb{K}^{n \times n}$

$$\|M\|_F^2 = \text{trace}(M^T M) = \sum_{i=1}^n \sigma_i(M^T M) = \sum_{i=1}^n \sigma_i^2(M)$$

because the eigenvalues and singular values of $M^T M$ coincide, so (6) implies the bounds in (4). Furthermore,

$$\|(I - E)^{-1}\|_F = \|I + \sum_{\nu=1}^{\infty} E^\nu\|_F \quad \text{and} \quad \left\| \sum_{\nu=1}^{\infty} E^\nu \right\|_F \leq \frac{\alpha}{1 - \alpha}$$

imply the bounds in (5). Using $\|E\|_2 \leq \|E\|_F$ implies

$$\frac{\text{left (5)}}{\text{left (4)}} \leq \frac{\sqrt{n} - \alpha/(1 - \alpha)}{\sqrt{n}/(1 + \alpha)} \leq 1 + \alpha < 2$$

and

$$\frac{\text{right (4)}}{\text{right (5)}} \leq \frac{\sqrt{n}}{\sqrt{n}(1 - \alpha) + \alpha} = \frac{1}{1 - \alpha(1 - 1/\sqrt{n})} \leq \sqrt{n} .$$

For α close to 1, the bounds in (5) tend to $\pm\infty$ whereas, with a gap between $\|E\|_2$ and $\|E\|_F$, the bounds in (4) remain finite. A simple example is again E being a multiple of the identity matrix. This completes the proof. \square

2.2 General Point Matrices

Bounds for the condition number of a point matrix $A \in \mathbb{K}^{n \times n}$ follow by direct application of the previous lemma.

Lemma 2.3 *Let $A, R \in \mathbb{K}^{n \times n}$ and $p \in \{1, 2, \infty, F\}$ be given, abbreviate $E := I - RA$ and assume $\|E\|_p \leq \alpha$ for some $\alpha \in \mathbb{R}$. Suppose $\alpha < 1$ in the cases $p \in \{1, 2, \infty\}$, and $\|E\|_2 < 1$ in the case $p = F$. Then A and R are non-singular and*

$$\frac{\|A\|_p \|R\|_p}{1 + \alpha} \leq \kappa_p(A) \leq \frac{\|A\|_p \|R\|_p}{1 - \alpha} \quad \text{for } p \in \{1, 2, \infty\}, \quad (7)$$

and

$$\|A\|_F \|R\|_F (1 - \alpha\beta) \leq \kappa_F(A) \leq \|A\|_F \|R\|_F (1 + \alpha\beta) \quad (8)$$

for $\beta := \sqrt{n}/(1 - \|E\|_2)$. If $\alpha < 1$ in the case $p = F$, then (8) is also true for $\beta = \sqrt{n} + \alpha/(1 - \alpha)$.

PROOF. The assumptions imply $I - E$ and therefore A and R are invertible. Using

$$A^{-1} = (I - E)^{-1}R = R + (I - E)^{-1}ER$$

together with Lemmas 2.1 and 2.2 yields the right bound in (7) and both bounds in (8), whereas the left bound in (7) follows by

$$\|R\|_p = \|(I - E)(I - E)^{-1}R\|_p = \|(I - E)A^{-1}\|_p \leq (1 + \alpha)\|A^{-1}\|_p.$$

□

Note that for both lower bounds, only upper bounds of $\|(I - E)^{-1}\|$ have been used. For the lower bound in (7), we could use the splitting of 1 since $\|I\|_p = 1$ for $p \in \{1, 2, \infty\}$; this is not possible for the lower bound in (8). One might be inclined to expand $(I - E)^{-1} = I + (I - E)^{-1}E$ to improve the bounds in (7) as well. However, this yields

$$\|R\| \left(1 - \frac{\alpha}{1 - \alpha}\right) \leq \|A^{-1}\| \leq \|R\| \left(1 + \frac{\alpha}{1 - \alpha}\right) \quad (9)$$

implying the same upper, but a weaker lower bound. This is due to the mentioned asymmetry of the bounds in Lemma 2.1.

2.3 General Interval Matrices

Let $\mathbf{A} \in \mathbb{IK}^{n \times n}$ be given, so that

$$\mathbf{A} = \langle M, \Delta \rangle := \{\tilde{A} \in \mathbb{K}^{n \times n} : |\tilde{A} - M| \leq \Delta\} \quad \text{for } M \in \mathbb{K}^{n \times n} \text{ and } 0 \leq \Delta \in \mathbb{R}^{n \times n}. \quad (10)$$

Let $\tilde{A} \in \mathbf{A}$ be fixed but arbitrary, so that $\tilde{A} = M + \tilde{\Delta}$ with $|\tilde{\Delta}| \leq \Delta$. Denoting $\tilde{E} := I - R\tilde{A}$ and using (2) it follows

$$\|\tilde{E}\| \leq \|I - RM\| + \|R\| \|\tilde{\Delta}\| \leq \|I - RM\| + \|R\| \|\Delta\|.$$

Furthermore, again using (2),

$$\|M\|_p - \|\Delta\|_p \leq \|M\|_p - \|\tilde{\Delta}\|_p \leq \|\tilde{A}\|_p \leq \|M\|_p + \|\tilde{\Delta}\|_p \leq \|M\|_p + \|\Delta\|_p.$$

Proceeding as before by using $\tilde{A}^{-1} = (I - \tilde{E})^{-1}R = R + (I - \tilde{E})^{-1}\tilde{E}R$, we obtain

Lemma 2.4 *Let \mathbf{A} be as in (10), $R \in \mathbb{K}^{n \times n}$ and $p \in \{1, 2, \infty, F\}$ be given, and assume $\|I - RM\|_p + \|R\|_p \|\Delta\|_p < \gamma_p$ for some $\gamma_p \in \mathbb{R}$. Suppose $\gamma_p < 1$ in case $p \in \{1, 2, \infty\}$, and $\gamma_2 < 1$ in case $p = F$. Then for every $\tilde{A} \in \mathbf{A}$*

$$\frac{(\|M\|_p - \|\Delta\|_p)\|R\|_p}{1 + \gamma_p} \leq \kappa_p(\tilde{A}) \leq \frac{(\|M\|_p + \|\Delta\|_p)\|R\|_p}{1 - \gamma_p} \quad \text{for } p \in \{1, 2, \infty\}, \quad (11)$$

and

$$\left(\|M\|_F - \|\Delta\|_F\right)\|R\|_F(1 - \gamma_F\beta) \leq \kappa_F(\tilde{A}) \leq \left(\|M\|_F + \|\Delta\|_F\right)\|R\|_F(1 + \gamma_F\beta) \quad (12)$$

for $\beta := \sqrt{n}/(1 - \gamma_2)$. If $\gamma_F < 1$, then (12) is also true for $\beta = \sqrt{n} + \gamma_F/(1 - \gamma_F)$.

Again one might try to improve the bounds using $\tilde{A}^{-1} = (I + M^{-1}\tilde{\Delta})^{-1}M^{-1}$ and the implicit estimate for $\|M^{-1}\|$ in (7). But again a computation yields that the upper bounds coincide, whereas the lower bound is weaker than in (11).

2.4 Extremely Ill-Conditioned Matrices

We finally address matrices with $\kappa(A) > \text{eps}^{-1}$, where eps denotes the relative rounding error unit. In IEEE 754 binary64 (double precision), this means $\kappa(A) > 10^{16}$. In that case there is likely no preconditioner $R \in \mathbb{K}^{n \times n}$ at all such that $I - RM$ is convergent.

Here a method as in [9, Section 6], which is based on unpublished results of the mid 1980's (see the citations in [9]), can be used. If $\kappa(A) > \text{eps}^{-1}$, then an approximate inverse has, in general, no correct digits, but it contains enough information to ensure $\kappa(RA) \approx \text{eps} \cdot \kappa(A)$. Denote $B := RA$, and let S be an approximate inverse of B . Abbreviating with $F := I - SB$ and assuming $\|F\|_p \leq \alpha < 1$, we obtain as before $A^{-1} = B^{-1}R = (I - F)^{-1}SR$, and Lemma 2.3 implies

$$\frac{\|A\|_p\|SR\|_p}{1 + \alpha} \leq \kappa_p(A) \leq \frac{\|A\|_p\|SR\|_p}{1 - \alpha} \quad (13)$$

for $p \in \{1, 2, \infty\}$, and similarly for the Frobenius norm using Lemma 2.2.

It is seen in the final section that this approach works up to condition numbers near eps^{-2} , which is up to 10^{32} in IEEE 754 binary64. Note that this is far beyond the scope of common numerical algorithms. It is mandatory in that case to use extra-precise accumulation of dot products for computing the error bounds for $B = RA$ (and only there). A number of such algorithms using only floating-point operations in the current working precision are available, see [8] and the literature cited therein.

We close this section with some hints on the implementation for $p = 2$, the spectral condition number. The spectral norm of a matrix can be bounded using the methods in [8], however, this requires $\mathcal{O}(n^3)$ operations. In contrast, bounds for the spectral norm of a non-negative matrix C are computed in $\mathcal{O}(n^2)$ operations using $\|C\|_2^2 = \rho(C^T C)$, $C^T C \geq 0$ and Perron-Frobenius Theory. The inequalities

$$\min_i \frac{(C^T C x)_i}{x_i} \leq \rho(C^T C) \leq \max_i \frac{(C^T C x)_i}{x_i} \quad (14)$$

are valid for any positive vector x , see [1]. Here $(\cdot)_i$ denotes the i -th entry of a vector. Since the vector x is arbitrary, a power set iteration replacing x by $C^T(Cx)$ may be performed. Each iteration requires $\mathcal{O}(n^2)$ floating-point operations. It is known

[10] that this produces a nested sequence of lower and upper bounds, guaranteed to improve in each step.

The components of the matrix $E = I - RA$ or $I - RM$ are, for not too ill-conditioned matrices, likely to be small, and usually it does no harm to bound the norm by $\| |E| \|$. To this end (14) can be used, and also for bounding $\| |\Delta| \|$. This method should not be used to bound $\|A\|$ or $\|R\|$.

Note that all approaches use an approximate inverse, so they may be time and memory consuming for sparse matrices. However, error bounds for the condition number of a matrix imply bounds for the solution of linear systems. For general, sparse matrices, this is one of the grand challenges in verification methods [6]. The presented algorithms will be included in Version 7.2 of INTLAB [7].

3 Computational Results

We now show some computational results. Matrices of specified condition number are generated by predefining a geometric sequence of singular values as in the Matlab function `randsvd`; for extremely ill-conditioned matrices, we use the INTLAB-function `randmat`, see [9].

n	condition number						fail
	10^2	10^5	10^{10}	10^{13}	10^{14}	10^{15}	
10	$3.1 \cdot 10^{-14}$	$1.6 \cdot 10^{-11}$	$1.5 \cdot 10^{-6}$	$1.5 \cdot 10^{-3}$	$1.4 \cdot 10^{-2}$	$1.3 \cdot 10^{-1}$	0
20	$5.9 \cdot 10^{-14}$	$2.7 \cdot 10^{-11}$	$2.1 \cdot 10^{-6}$	$2.0 \cdot 10^{-3}$	$1.9 \cdot 10^{-2}$	$1.9 \cdot 10^{-1}$	0
50	$1.9 \cdot 10^{-13}$	$6.6 \cdot 10^{-11}$	$4.3 \cdot 10^{-6}$	$3.5 \cdot 10^{-3}$	$3.9 \cdot 10^{-2}$	$4.2 \cdot 10^{-1}$	0
100	$1.0 \cdot 10^{-5}$	$1.0 \cdot 10^{-5}$	$1.9 \cdot 10^{-5}$	$7.8 \cdot 10^{-3}$	$7.1 \cdot 10^{-2}$	1.2	5
200	$1.0 \cdot 10^{-5}$	$1.0 \cdot 10^{-5}$	$3.1 \cdot 10^{-5}$	$1.6 \cdot 10^{-2}$	$1.6 \cdot 10^{-1}$	30	95
500	$1.0 \cdot 10^{-5}$	$1.0 \cdot 10^{-5}$	$6.5 \cdot 10^{-5}$	$4.1 \cdot 10^{-2}$	$4.6 \cdot 10^{-1}$	-	100
1000	$1.0 \cdot 10^{-5}$	$1.0 \cdot 10^{-5}$	$9.7 \cdot 10^{-5}$	$6.5 \cdot 10^{-2}$	$9.4 \cdot 10^{-1}$	-	100

Table 1: Test for $p = 2$, median over 100 samples, number of failures for condition number 10^{15} in the last column.

n	condition number						fail	fail
	10^2	10^5	10^{10}	10^{13}	10^{14}	10^{15}		
10	$3.8 \cdot 10^{-14}$	$2.4 \cdot 10^{-11}$	$2.3 \cdot 10^{-6}$	$2.2 \cdot 10^{-3}$	$2.1 \cdot 10^{-2}$	0	$2.1 \cdot 10^{-1}$	0
20	$7.4 \cdot 10^{-14}$	$4.5 \cdot 10^{-11}$	$3.6 \cdot 10^{-6}$	$3.5 \cdot 10^{-3}$	$3.3 \cdot 10^{-2}$	0	$3.7 \cdot 10^{-1}$	0
50	$2.2 \cdot 10^{-13}$	$1.2 \cdot 10^{-10}$	$8.3 \cdot 10^{-6}$	$7.0 \cdot 10^{-3}$	$7.0 \cdot 10^{-2}$	0	1.1	11
100	$5.4 \cdot 10^{-13}$	$2.6 \cdot 10^{-10}$	$1.6 \cdot 10^{-5}$	$1.4 \cdot 10^{-2}$	$1.4 \cdot 10^{-1}$	0	3.6	87
200	$1.4 \cdot 10^{-12}$	$6.2 \cdot 10^{-10}$	$3.7 \cdot 10^{-5}$	$3.1 \cdot 10^{-2}$	$3.3 \cdot 10^{-1}$	0	-	100
500	$9.2 \cdot 10^{-12}$	$3.9 \cdot 10^{-9}$	$2.3 \cdot 10^{-4}$	$2.0 \cdot 10^{-1}$	-	100	-	100
1000	$2.1 \cdot 10^{-11}$	$7.7 \cdot 10^{-9}$	$3.7 \cdot 10^{-4}$	$3.4 \cdot 10^{-1}$	-	100	-	100

Table 2: Test for $p = 1$, median over 100 samples, number of failures for condition numbers 10^{14} and 10^{15} .

For a given matrix A and computed bound $\mathbf{c} \in \mathbb{IR}$ for $\kappa(A)$ the following tables show the relative error $\phi(\mathbf{c}, c) := \text{rad}(\mathbf{c})/c$, where $\text{rad}(\mathbf{c}) := (\bar{c} - \underline{c})/2$ denotes the radius of $\mathbf{c} = [\underline{c}, \bar{c}]$. Here $c = \text{cond}(A)$ is the approximate condition number computed by

ordinary Matlab for anticipated condition numbers up to 10^{15} , and using the symbolic toolbox for higher condition numbers. The quality of the approximation seems still acceptable for $\kappa(A) \lesssim 10^{15}$.

Tables 1 and 2 show the results for the spectral norm and the 1-norm, respectively. Bounds for the condition number are calculated for 100 samples, and the median of $\phi(\mathbf{c}, c)$ of the bounds is displayed. The lower bound of \mathbf{c} is set to zero if it becomes negative.

For the spectral norm, the method works in all examples up to condition number 10^{14} and, for $p = 1$, always up to condition number 10^{13} in our examples. If there are failures for a certain condition number, the number of failures is displayed in the following column. For the spectral norm and dimensions up to 50, bounds for the norm are computed on the base of a full eigendecomposition, otherwise the accuracy is restricted to about 10^{-6} to save computing time. This is because an inclusion of the condition accurate to a few digits seems appropriate and sufficient in almost all applications.

Despite this the computational results are consistent with the fact that the condition number of the condition number is the condition number [2]. The results for the Frobenius norm are a little bit better in accuracy than those for the 1-norm, and (of course) equal to those for the spectral norm in terms of failures.

Tables 3 and 4 show the results for the spectral norm and the 1-norm, respectively, for extremely ill-conditioned matrices. As has been mentioned, (only) the product RA is computed with extra-precise dot product accumulation. It means that each entry $(RA)_{ij}$ is computed with algorithm `Dot2` in [8]. This algorithm uses only floating-point operations in working precision and produces a result *as if* computed in twice the working precision. Otherwise the bounds in (13) are straightforwardly computed.

The results are as expected in terms of accuracy and failure. As before, the results for the Frobenius norm are similar in accuracy to those of the 1-norm, and equal to those the spectral norm in terms of failure.

n	condition number						
	10^{15}	10^{20}	10^{25}	10^{28}	fail	10^{29}	fail
10	$7.2 \cdot 10^{-15}$	$1.7 \cdot 10^{-13}$	$9.3 \cdot 10^{-9}$	$6.7 \cdot 10^{-6}$	0	$3.3 \cdot 10^{-6}$	0
20	$1.3 \cdot 10^{-14}$	$2.1 \cdot 10^{-13}$	$3.7 \cdot 10^{-8}$	$1.6 \cdot 10^{-5}$	1	$3.0 \cdot 10^{-5}$	0
50	$3.5 \cdot 10^{-14}$	$1.3 \cdot 10^{-12}$	$3.0 \cdot 10^{-8}$	$1.6 \cdot 10^{-4}$	1	$2.6 \cdot 10^{-1}$	91
100	$1.0 \cdot 10^{-5}$	$1.5 \cdot 10^{-5}$	$2.0 \cdot 10^{-4}$	$8.5 \cdot 10^{-2}$	26	$8.8 \cdot 10^{-1}$	99

Table 3: Test for $p = 2$, median over 100 samples, number of failures for condition numbers 10^{28} and 10^{29} .

n	condition number						
	10^{15}	10^{20}	10^{25}	10^{28}	fail	10^{29}	fail
10	$3.1 \cdot 10^{-15}$	$1.9 \cdot 10^{-13}$	$7.1 \cdot 10^{-9}$	$5.1 \cdot 10^{-6}$	0	$3.6 \cdot 10^{-6}$	0
20	$5.4 \cdot 10^{-14}$	$2.3 \cdot 10^{-13}$	$3.2 \cdot 10^{-8}$	$1.1 \cdot 10^{-5}$	1	$1.5 \cdot 10^{-5}$	1
50	$1.2 \cdot 10^{-14}$	$1.3 \cdot 10^{-12}$	$3.9 \cdot 10^{-8}$	$1.8 \cdot 10^{-4}$	0	$3.0 \cdot 10^{-1}$	92
100	$2.4 \cdot 10^{-14}$	$1.4 \cdot 10^{-9}$	$1.6 \cdot 10^{-4}$	$7.1 \cdot 10^{-2}$	15	$8.2 \cdot 10^{-1}$	97

Table 4: Test for $p = 1$, median over 100 samples, number of failures for condition numbers 10^{28} and 10^{29} .

We finally mention that using (1) produces, in general, more accurate inclusions but needs more computing time. Table 5 shows the median of the computing times of 100 examples each. We choose some random matrices because there is not much difference in time for well- or ill-conditioned matrices. As can be seen, the new approach is faster than (1). The timing for interval matrices is very similar and therefore omitted. Matrices with condition number beyond eps^{-1} are outside the scope of (1).

	$n = 100$		$n = 200$		$n = 500$		$n = 1000$	
	new	(1)	new	(1)	new	(1)	new	(1)
$\ \cdot\ _1$	0.0034	0.0089	0.0092	0.032	0.070	0.35	0.35	2.01
$\ \cdot\ _2$	0.0141	0.0196	0.0276	0.054	0.158	0.45	0.93	2.57
$\ \cdot\ _\infty$	0.0038	0.0092	0.0090	0.035	0.070	0.35	0.35	2.00
$\ \cdot\ _F$	0.0051	0.0098	0.0089	0.033	0.065	0.34	0.34	1.97

Table 5: Computing time in seconds new vs. (1)

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